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None

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(54) Object lifetime control in computer memory protection.

(57) An object-oriented computer architecture in which access descriptors (24) include an object index (34) for selecting an object in the address space, and a rights field (35) specifying the permissible operations on a bi-paged object (38) selected by the access descriptor. A local object lifetime bit is provided in the encoded fields portion (35, 50, 62, 81) of access descriptors, object descriptors, and page table entries to determine the lifetime of an object. The AD lifetime bit in the encoded fields (35) of AD (24) is compared in OTE Lifetime Check Logic with the destination object lifetime, the OTE local bit in the encoded fields (50) of the OTE access descriptor. The OTE local bit in the encoded fields (50) of the OTE is compared in PDTE Lifetime Check Logic with the destination object lifetime, the PDTE local bit in the encoded fields (62) of the PDTE access descriptor. The PDTE local bit in the encoded fields (62) of the PDTE is compared in PTE Lifetime Check Logic with the destination object lifetime, the PTE local bit in the encoded fields (81) of the PTE access descriptor. If any of these checks fails, a protection lifetime fault is asserted.

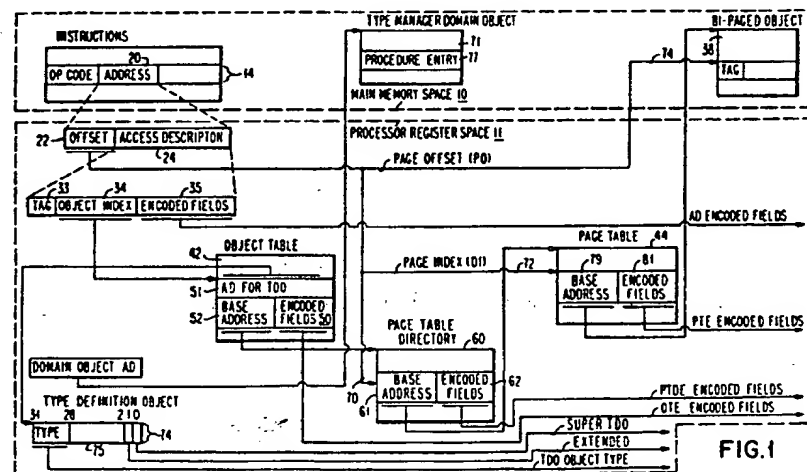
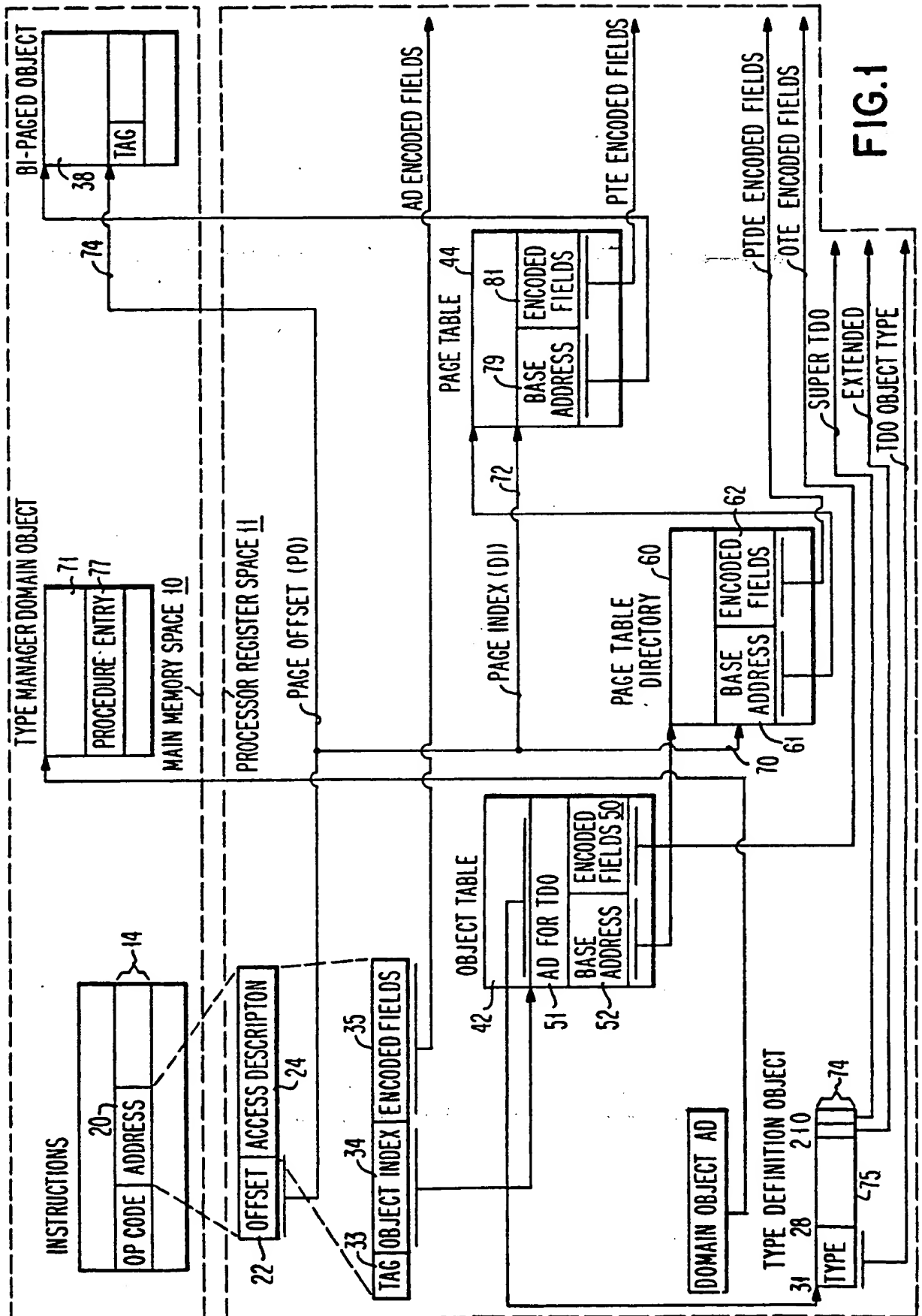


FIG. 1



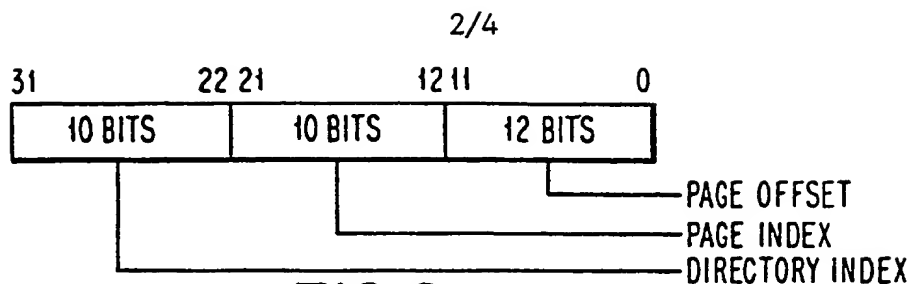


FIG. 2

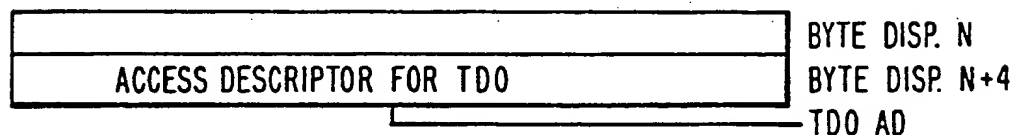


FIG. 3

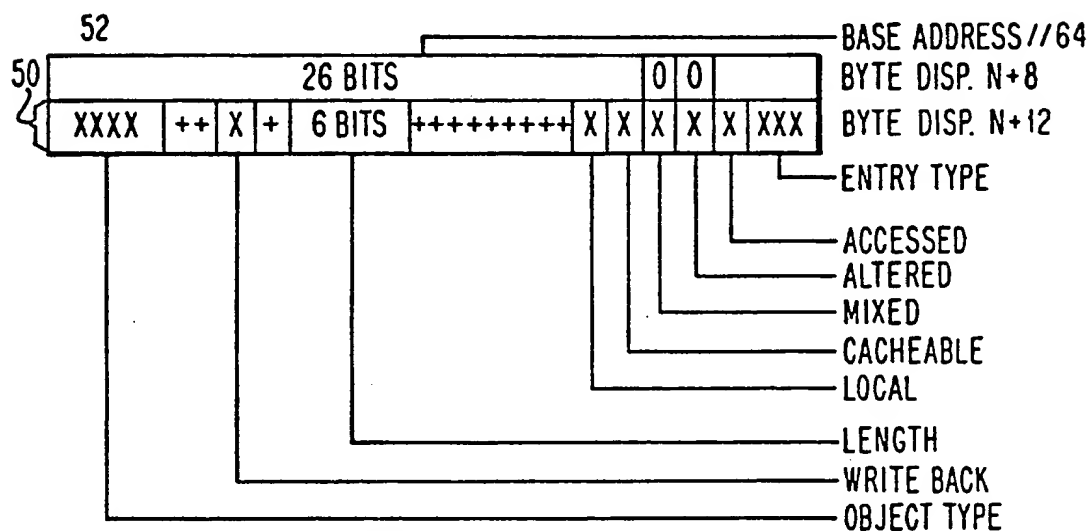


FIG. 4

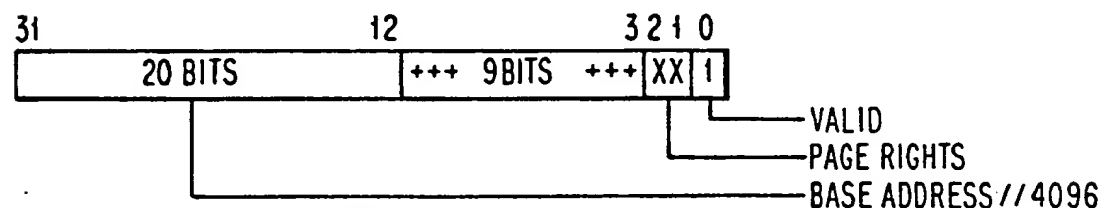


FIG. 5

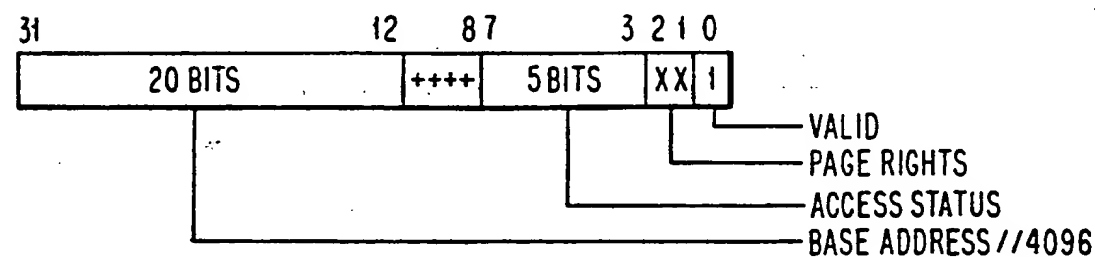


FIG. 6

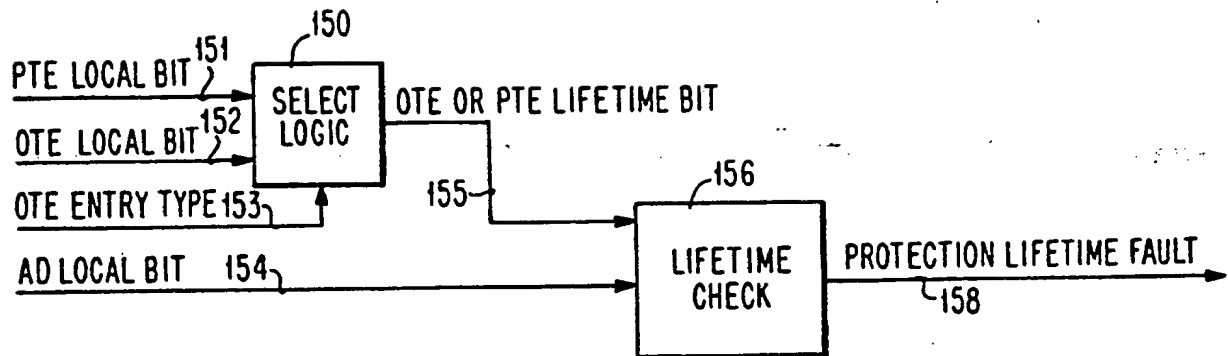


FIG. 7

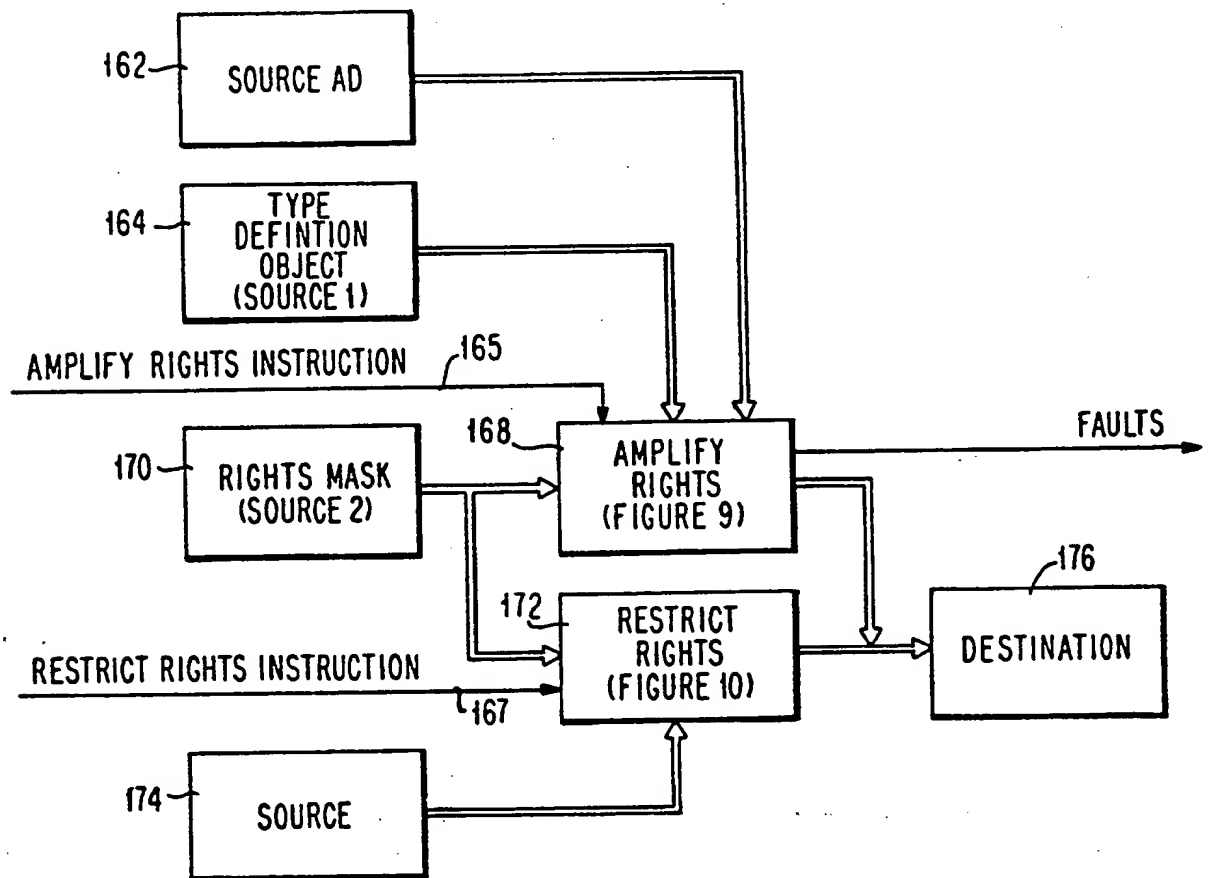


FIG.8

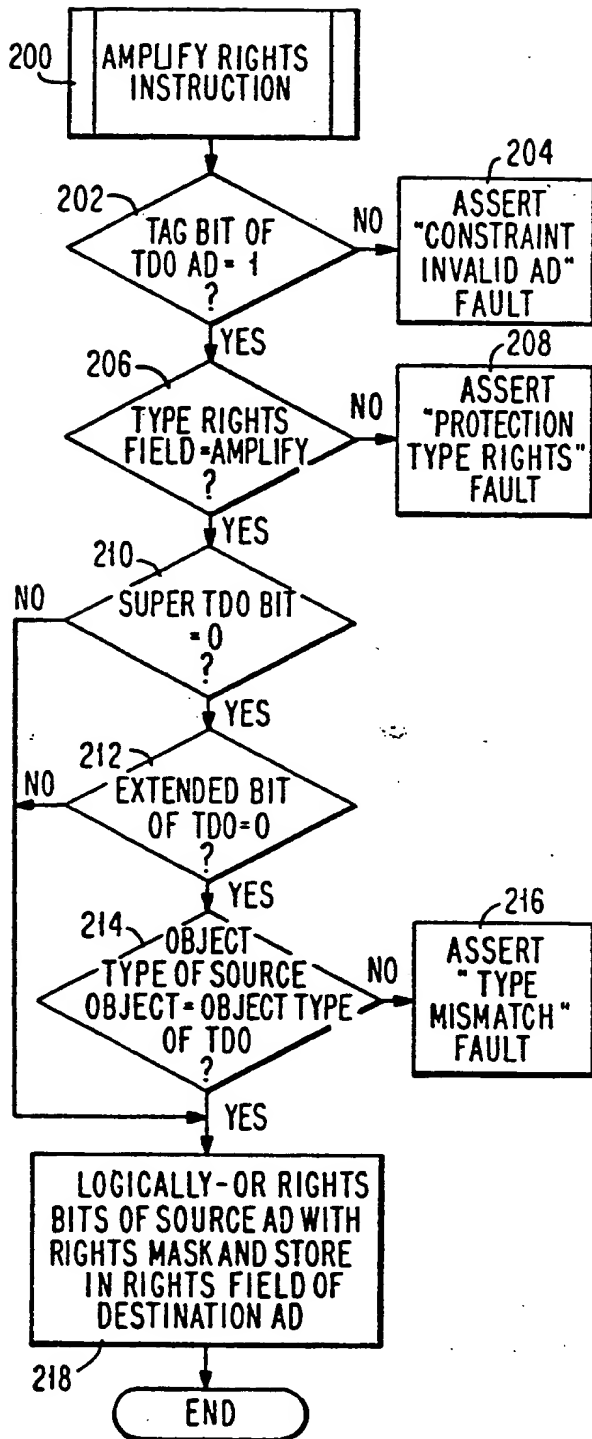


FIG. 9

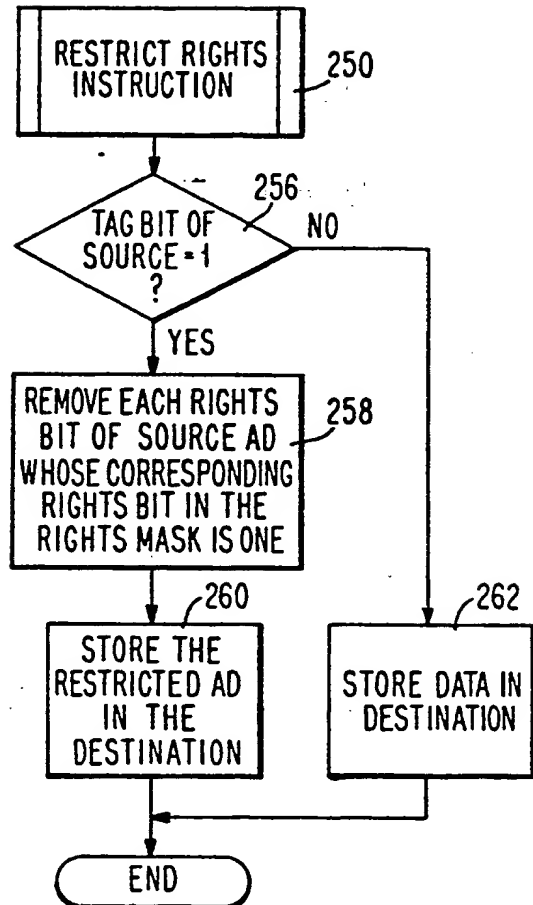


FIG. 10

OBJECT LIFETIME CONTROL IN AN OBJECT-ORIENTED MEMORY PROTECTION MECHANISM

CROSS REFERENCE TO RELATED APPLICATIONS

Application SN 07/455,625 (D-1256) filed 12/22/89 entitled "Disabling Tag Bit Recognition and Allowing Privileged Operations to Occur in an Object-Oriented Memory Protection Mechanism", assigned to Intel Corporation, the assignee of the present invention.

Application SN 07/455,626 (D-1258) filed 12/22/89 entitled "Synchronous Communication Between Execution Environments in a Data Processing System Employing an Object-Oriented Memory Protection Mechanism", assigned to Intel Corporation.

Application SN 07/455,635 (D-1259) filed 12/22/89 entitled "Type Management and Control in an Object-Oriented Memory Protection Mechanism", assigned to Intel Corporation, the assignee of the present invention.

BACKGROUND OF THE INVENTION

Field of the Invention

The invention relates to data processing systems, and more particularly to a memory-accessing system which deallocates objects stored in memory in accordance with assigned object lifetimes.

Description of the Related Art

Modern computer systems are organized such that all information within the system is represented by hardware-recognized, memory-resident information structures called objects. This is known as an object-oriented computer architecture. An object is a package of data that is a represen-

tation of related information maintained in a contiguously-addressed set of memory locations.

In an object-oriented programming system, each type or class of object has a set of operations that are permitted to be performed on the objects that are members of the class. This is known as an object-based access mechanism. The internal structure of objects is hidden from the programmer. The programmer can manipulate the objects only with the permitted operations. Access to an object is by means of a pointer to the object called an access descriptor. An access descriptor describes the kind of representation (rep) rights (such as read rights or write rights) which define the various kinds of access to the object that are permitted by the holder of the access descriptor for that object.

The above described mechanisms allow many processes running on the system to have access to a very large shared data space. When a process terminates, objects stored in the shared data space need to be reclaimed (made available for use again). This is usually done at great expense in terms of operating system overhead by a process known as "garbage collection". This involves perhaps an additional processor, operating concurrently with the main processor, which detects and collects garbage. It must identify those objects whose content is no longer useful for the task in progress and make the associated memory locations available again for some other use. This can be very time consuming and therefore expensive in term of processing overhead and reduced throughput.

It is an object of this invention to provide a mechanism for identifying those objects associated with a process whose content is no longer useful and make the associated memory locations available again for some other use when the process terminates.

It is a further object of this invention to provide a mechanism for assigning a lifetime to objects associated with a process to thereby enable the implicit deallocation of certain objects and to prevent dangling references if access descriptors which point to objects with shorter lifetimes are stored in objects with longer lifetimes.

SUMMARY OF THE INVENTION

Briefly, the above objects are accomplished in accordance with the invention by providing an object table object (42), a page table directory object (60), and a page table object (44). Each of the objects is comprised of words, the words being either access descriptors (30) or data words (32). Each access descriptor includes an object index for selecting an object in the address space of the processor's memory, and a rights field specifying the permissible operations on an object selected by the access descriptor. The object table object (42) is stored at an address accessible by the processor upon processor initialization and has stored therein object descriptors for use by the processor in forming physical addresses to the page table directory object (60). Each of the access descriptors has a lifetime bit indicating the lifetime of an object specified by the access descriptor and the object storage descriptor has a lifetime bit indicating the lifetime of the object selected by the access descriptor.

SUMMARY OF THE INVENTION

Briefly, the above objects are accomplished in accordance with the invention by providing an object table object (42), a page table directory object (60), and a page table object (44). Each of the objects is comprised of words, the words being either access descriptors or data words. Each access descriptor includes an object index for select-

ing an object in an the address space of the processor's memory, and a rights field specifying the permissible operations on an object selected by the access descriptor. The object table object (42) is stored at an address accessible by the processor upon processor initialization and has stored therein object descriptors for use by the processor in forming physical addresses to the page table directory object (60). Each of the object descriptors includes a first encoded field (52) specifying the base address of the object referenced by the object descriptor and a second encoded field (50) which includes object entry type. The page table (44) has stored therein page table entries for use by the processor in forming physical addresses to the paged object (38) and each page table entry has a page rights field (81). The offset (22) includes a page index (PI) used to index into a selected page table to locate the page table entry, and a page offset (PO) used in conjunction with a base address (79) stored in the page table entry to locate a datum entry in the paged object (38). Each of the access descriptors has a lifetime bit indicating the lifetime of an object specified by the access descriptor and the object storage descriptor or page table entry has a lifetime bit indicating the lifetime of the object or page selected by the access descriptor and offset.

The invention has the advantage that objects can be reclaimed without resort to the expensive process of garbage collection. With the support of the local lifetime, it is possible to reclaim all local objects associated with a specific local execution environment (process) without the overhead of garbage collection when the local environment terminates.

Brief Description of the Drawings

The foregoing and other objects, features, and advantages of the invention will be apparent from the following detailed description of a preferred embodiment of the invention, as illustrated in the accompanying drawings wherein:

FIGURE 1 is a block diagram of the processor registers and of system objects for supporting object addressing and protection;

FIGURE 2 is a more detailed block diagram of the object offset of FIGURE 1;

FIGURE 3 and FIGURE 4 together comprise a diagram of a storage descriptor within the object table of FIGURE 1;

FIGURE 5 is a diagram of the page table directory entry (PTDE) within the page table directory FIGURE 1; and,

FIGURE 6 is a diagram of the page table entry (PTE) within the page table FIGURE 1;

FIGURE 7 is a diagram of lifetime protection checking logic;

FIGURE 8 is a diagram of the execution logic of an amplify_rights instruction and a restrict_rights instruction;

FIGURE 9 is a diagram of data flow within the amplify rights logic of FIGURE 8; and,

FIGURE 10 is a diagram of data flow within the restrict rights logic of FIGURE 8.

Description of the Preferred Embodiment

The microprocessor in which the present invention may be embodied is shown in more detail in patent application SN 07/455,625 D-1256). To summarize briefly, the microprocessor is logically subdivided into the following major units: an Instruction Fetch Unit, an Instruction Decoder and Microinstruction Sequencer, a Translation Lookaside Buffer, a Local Bus Sequencer, and an Integer Execution Unit and Floating Point Unit.

The Translation Lookaside Buffer (TLB) performs the address translation needed to implement virtual memory mechanisms. The address translation and memory protection uses an associative table of storage descriptors and page table entries. The TLB contains an address cache, an address adder, and memory protection checking hardware.

Execution Modes

As described in patent application SN 07/455,625 (D-1256), the system provides two modes of execution, User mode and Supervisor mode, to support the efficient emulation of conventional operating systems. The system does not define the notion of privileged instructions. All instructions can be executed in either mode. A program gains privilege by nature of its access representation (rep) rights and its execution mode. The page rep rights in the current linear address space are interpreted differently depending on the execution mode. Storage used by the operating system generally has page rep rights which do not allow user access, but may be read-only or read/write in the Supervisor mode.

An instruction (14) in the memory may contain an address (20) which is a virtual memory address to the object (38). The address is stored in a processor register (11) in the

processor. The address includes an offset (22) and an access descriptor (24). The access descriptor (24) is a 32-bit word which has a tag bit (33). The access descriptor includes an object index (34) and a rights field (35). The object index provides an offset into an object table (42) which contains storage descriptors for objects stored in the memory, one of the objects being a page table (44). The descriptor includes an encoded field (50) and a base address (52). An object is described by an object table entry in the object table (42) which provides the mapping information for the physical addresses of the storage blocks and pages which together represent the physical objects. For a simple object, the object is represented by a block in physical address space directly. The physical base address (52) is stored in the object table entry. The information within the object is located by using a simple offset in the virtual address.

A paged object is represented by a set of physical pages using a single-level page table. A bi-paged object (38) is represented by a set of physical pages using a page table (44) and a page table directory (60). The object table entry for a paged object, called a paged object descriptor, contains the physical base address (61) of the page table directory (60), which is an array of page table directory entries for the page tables, including page table (44). The page table directory entry is found by indexing using the page table directory index (DI) located in the offset (22). The page table entry is found by indexing using the page index (PI) located in the offset (22). The bi-paged object entry is found by indexing using the page offset (PO) located in the offset (22).

Object Addressing

Three different address spaces are defined, Logical Address Space, Virtual Address Space and Physical Address Space. The Logical Address Space is accessed with a Linear Address which includes an Object Selector and an Object Offset. The Virtual Address Space is accessed with an Object Index and an Object Offset. The Physical Address Space is accessed with an Object/Page Table Entry and Block/Page Offset.

A logical address space is mapped onto a virtual address space which is mapped onto a physical address space. An address in each space has a unique structure as described below.

Physical Address Space

The physical address space covers read-write memory, read-only memory, and memory-mapped I/O. Communication with the physical world is in terms of physical addresses. The physical address space is byte addressable and must guarantee atomic and indivisible access (read or write) for memory addresses that are aligned according to the size of the access, i.e. a word access must be aligned to a word boundary. An indivisible access guarantees that a processor reading or writing a set of memory locations will complete the operation before another processor can read or write the same location. An atomic operation allows a processor to read and modify a set of memory locations with the guarantee that another processor doing an atomic operation on the same block will be delayed.

Virtual Address Space

The virtual address space is a collection of independent address spaces called objects. Any datum within an object is

located by a simple offset. The virtual address of the datum is specified by two components: an object index that selects the desired address space and an object offset that selects the datum within the address space. The size of the virtual address space is the product of the number of objects allowed and the maximum size of each object.

An object is also defined as the unit of protection. To control access within the virtual address space, the system protects and restricts the generation of object indices. An access descriptor (AD), is defined to contain an object index. The system protects the AD pointer by not allowing ordinary data to be used to represent it. An access descriptor (AD) can point to any of the objects in the virtual address space. The virtual address is, therefore, specified by a protected object index (i.e. an access descriptor) and an unprotected offset into the object.

Logical Address Space

Access descriptors, directly or indirectly accessible, are conceptually assembled in sets to form yet a third type of address space called a logical (or instantaneous) address space. The logical address space defines the visibility of the execution environment. Logical addresses are mapped onto the single virtual address space. For maximum flexibility, two types of logical addresses are defined: linear address and structured address.

The first type, the linear address, is defined by four objects that form the execution environment. A linear address is used to represent the conventional notion of a process address space. Linear addresses, interpreted within a given environment, are mapped onto the virtual address space. The mapping of linear addresses to virtual addresses is a fundamental part of the instruction interpretation process. In

a linear address, an operand specifier supplies only an offset; the current linear address space is implied. The upper two bits of a linear address implicitly selects one of the four objects that define the execution environment, while the remaining 30 bits is an offset into the selected object.

The second type, the structured address, is defined by a virtual address (ie. access descriptor plus offset). The structured address is used to invoke advanced object-oriented protection features. In a structured address, an operand specifier supplies a virtual address. Since the system does not allow specifying an AD directly in the instruction stream, the AD part of the virtual address must be specified indirectly using an AD selector in an operand specifier within the instruction. An AD selector specifies an AD stored locally in the processor registers. This limits the instantaneous addressability of objects to those accessible by the ADs in the processor registers only.

Access Descriptors and their Rights

An access descriptor (AD) is a protected pointer to an object in the address space. The system protects access descriptors from accidental or malicious creation and modification.

A program cannot address an object directly, but only indirectly via an access descriptor in the execution environment. Since a program cannot reference an object without an access descriptor to it, a program's visibility can be restricted to only those objects it needs to access.

An access descriptor contains the following information: Object Index, Rights, and Lifetime bit. The Object Index selects the object. The rights specify read rights, write

rights and type rights in an AD. The rights indicate the permissible operations on the object. Rights are associated with an access descriptor and not with the object itself. It is thus possible to have different rights to the same object by selecting different access descriptors. The Lifetime bit indicates the lifetime of the object that an AD references.

AD to Object Mapping

Objects are referenced using access descriptors which are system-wide protected pointers. The object index (34) in an access descriptor selects an object table entry (OTE) in the system wide object table object (42). An object table entry specifies the location, size, type, etc. of the referenced object.

Storage Blocks and Pages

An object is physically composed of a storage block and/or a set of pages. A block is a contiguous area in the physical address space. A block can be used to represent a simple object, a page table, or a page table directory.

The base address of a storage block points to the first visible byte of the block. The base address of a storage block must be aligned on a 64-byte physical address boundary. The length of a block varies from 64 bytes to 4096 bytes. A block cannot span across a 4K byte boundary.

An object can also be represented by a set of pages with one or two level(s) of page tables. The first level table can be a storage block instead of a page. The pages that define the bi-paged object (38) are described by the page table (44). A page is a fixed size block of 4K bytes with base address aligned on a 4K-byte boundary.

Tagging

An object may contain access descriptors and/or data, i.e. any binary information. Access descriptors and data can reside in the same object and can be interleaved in any arbitrary order. In some systems, a tag bit is associated with each 4-byte aligned word in memory to indicate whether the word is data or possibly a valid access descriptor. A valid access descriptor must be aligned to a 4-byte boundary with its tag bit set to one. A tag bit of zero is used to indicate that the word contains data or an access descriptor that cannot be dereferenced, that is cannot be used to reference the contents of an object.

In some systems, the tag bit is not available. In such systems, the interpretation of a word as data or an access descriptor depends upon the operation.

In a word-aligned read or write of the whole word, the tag bit is either preserved or set to zero depending on the operation. In a non-word aligned read, or a partial read of a word, the tag bit of the returned value is always forced to zero. The data manipulation (arithmetic or logical) instructions generate results with tag bits forced to zero, independently of the tag bits of the source operands. Tagging is described more fully in patent application SN 07/455,625 (D-1256).

Typed Objects

The internal organization of certain objects is defined by specifying and interpreting data structures within their address spaces. These objects play a key role in the protection system, the interprocess/interprocessor communication system, and the storage management system. To recognize

these system-defined objects and to control their use, each one may be identified by a code. The code is called the object type and it is maintained with the object's address mapping information. The present invention allows users to invent their own object types and to assign them unique type codes.

Object Offset

An object offset is a 32-bit ordinal used to specify a datum within an object. It is capable of pointing to either data or access descriptors in an object. An object offset is divided into a number of fields. The interpretation of these fields is dependent on the object representation.

Object Size

The size of an object is specified in encoded fields (50) of the object table entry (OTE). The object offset (22) in a virtual address (20) plus the operand size is compared with the size of the referenced object on every address translation. This operation is called bounds checking and prevents reference beyond the specified object of a datum which may belong to another object.

Object Representation

An object is described by the object table entry which provides the mapping information for the physical addresses of the storage blocks and pages which together, directly or indirectly, represent the physical object. The system uses three different mapping schemes for different maximum object sizes and to minimize object representation overheads.

Simpl Objects. A simple object is represented by a block in physical address space directly. The physical base address

is stored directly in the object table entry. Such an entry is called a simple object descriptor.

Paged Objects. A paged object is represented by a set of physical pages using a single-level page table. The object table entry for a paged object, called a paged object descriptor, contains the physical address of a page table, which is an array of page table entries for the pages.

Bipaged Objects. A bipaged object is represented by a set of physical pages using two levels of page tables. As shown in FIGURE 1, the object table entry for a bipaged object, called a bipaged object descriptor, contains the physical address of a page table directory (60), which is an array of page table entries for page tables, including the page table (44) shown.

A simple object is defined by a simple object descriptor and represented by a single block. A simple object descriptor contains the physical base address and the block length. A simple object cannot span across a 4K byte physical address boundary.

A simple object offset is partitioned as follows:

Directory Index DI (bits 22-31). This 10-bit field must be zero. Otherwise a PROTECTION fault is raised.

Page Index PI (bits 12-21). This 10-bit field must be zero. Otherwise a PROTECTION fault is raised.

Block Offset SO (bits 0-11). This 12-bit field is the byte displacement added to the base address of the block to form the physical address for the first byte of the operand.

Paged Objects

A paged object is described by an object table entry called a paged-object descriptor. Paged objects are implemented with one level of page table. Each page is individually swappable and relocatable, thus not all pages of a paged object need be present in physical address space at the same time. To access an item of a paged object, only the page table and the selected page need to be located in the physical address space.

A paged-object descriptor contains the object length, but does not contain the base addresses (45) of the pages which represent the object. The base address field (45) of a paged-object descriptor contains the base address of the page table block. The length of the page table block is defined by the object length of the object.

A paged object offset (22) is partitioned as follows:

Directory Index, DI (bits 22-31). This 10-bit field must be zero. Otherwise, a PROTECTION fault is raised.

Page Index, PI (bits 12-21). This 10-bit field is used to index into the selected page table (44) for a page table entry.

Page Offset, PO (bits 0-11). This 12-bit field is the byte displacement and is appended to the base address (45) in the page table entry of the page to form the physical address for the first byte of the operand in the paged object (38).

Bipaged Objects

Refer to FIGURE 1. A bipaged object (38) is described by an object table entry (50, 52) called a bipaged object descrip-

tor. Bipaged objects are implemented with two levels of page tables. Each page or page table is individually swappable and relocatable, thus not all pages or page tables of a bipaged object need be present in physical address space at the same time. To access an item of a bipaged object, only the page table, the selected page table, and the selected page need to be located in the physical address space.

A bipaged object descriptor contains the object length, but does not contain the base addresses of the pages nor page tables which represent the object. The base address field (52) of a bipaged object descriptor contains the base address of the page table directory block (60). The length of the page table directory block is defined by the object length of the object.

A bipaged object offset is partitioned as follows:

Directory Index, DI (bits 22-31). The directory index (70) selects a page table entry in the page table directory specified by the bipaged object descriptor.

Page Index, PI (bits 12-21). The page index (72) selects a page table entry in the specified page table.

Page Offset, PO (bits 11-0). The page offset (74) is used as an offset into the page. The page offset is appended to the base address (in a page table entry) to form the physical address for the first byte of the operand.

Object Lifetime

To support the implicit deallocation of certain objects and to prevent dangling references, the system defines the concept of object lifetime. Dangling references can only occur if access descriptors which point to objects with shorter

lifetimes are stored in objects with longer lifetimes. The local object lifetime bit in the encoded fields portion of access descriptors, object descriptors, and page table entries are the means by which the system determines the lifetime of an object and prevents a potential dangling-reference access descriptor from being stored.

Local Bits

A local bit is associated with each object to denote its relative lifetime. The local bit is located in the encoded fields portion of the object descriptor for the object, for example in the encoded fields (50) of the storage descriptor in the object table (42). A value of 0 indicates a global object with unbound object lifetime. A value of 1 indicates a local object with bound object lifetime.

Lifetime Checking

Refer to FIGURE 7. The object lifetime check is performed every time an AD is stored. Since this requires the lifetime of the source AD and destination location to be compared, the operation is called lifetime checking. If a valid source AD has its associated local bit set to 1 and the lifetime of the destination location (in an object table entry or a page table entry) has its associated local bit set to 0 (is of global lifetime), a PROTECTION.LIFETIME fault is signalled. The destination location(s) may be modified if there is no other protection fault, but the AD of local lifetime must never be stored in the destination.

In FIGURE 7, the AD local bit in the encoded fields (35) of AD (24) is compared in OTE Lifetime Check Logic (150) with the destination object lifetime, the OTE local bit in the encoded fields (50) of the OTE access descriptor. The OTE local bit in the encoded fields (50) of the OTE is compared

in PDTE Lifetime Check Logic (152) with the destination object lifetime, the PDTE local bit in the encoded fields (62) of the PDTE access descriptor. The PDTE local bit in the encoded fields (62) of the PDTE is compared in PTE Lifetime Check Logic (154) with the destination object lifetime, the PTE local bit in the encoded fields (62) of the PTE access descriptor. If any of these checks fails, a protection lifetime fault (158) occurs at the output of the OR (156).

In the implicit manipulation of system objects, lifetime checking is ignored unless explicitly specified.

When an access descriptor is stored into memory, lifetime checking is performed first. When an access descriptor is written into memory, lifetime checking is not performed.

Mapping Tables

Object Table Objects

An object table object (42) serves as the root of the virtual address mapping. An object table is an array of 16-byte object table entries. The object index field (34) in the access descriptor (24) selects an object table entry in the object table (42). Object table entries are described subsequently.

Object tables do not have a system defined system type. Although an AD that is a pointer to an object table OT has global lifetime, the object table entries (OTES) or page table entries of an object table (OT) must have local lifetime. This is necessary to support Type Definition Object (TDO) ADs in OTEs and embedded descriptors.

There is one system wide object table (OT) for all processors which share a single system wide virtual address space.

System Defined Object Indices

The following object indices are defined:

<u>Object Indices</u>	<u>Purpose</u>
0	Reserved for empty AD.
1-7	Preserved
8	Object Table
9-15	Preserved
16-31	Default TDO ADs for object types 0-15

Page Tables or Page Table Directories

Page tables (44) are used to provide one or two level(s) of mapping for paged objects and bipaged objects. Page table directories (60) are used to provide the first level of mapping for bipaged objects. Page tables (or page table directories) contain page table entries (or page table directory entries) which define the base address of a page, and other information for virtual memory management and garbage collection.

Page tables and page table directories are system recognized, but are not objects and do not have a system type.

A page table (44) is an array of page table entries, each of which is 4 bytes in length. Each page table entry in a page table describes a page in a paged object or a bipaged object. Each page table entry in a page table directory (60) describes a page table (44) for a bipaged object (38).

The page table of a paged object or the page table directory of a bipaged object can be variable in size and aligned on any 64-byte boundary. The page tables of a bipaged object must be 4K bytes in size and aligned on 4K bytes boundary.

Page tables and page table directory are not objects and thus cannot be accessed directly in the virtual address space. In one embodiment of the invention they are accessed using physical addresses. In another embodiment of the invention the page tables are mapped to part of the object they are defining. In the second approach, the physical address of the page table directory or the page table is duplicated. Software guarantees that the physical address alias is updated during page swapping.

Descriptor Formats

Data Words

The fields of a data word are defined as follows:

Data (bits 0-31). This field contains any data value.

Tag (Tag Bit). This bit is 0 for data values.

Access Descriptors Access descriptor

The diagram of FIGURE 3 includes an access descriptor. The fields of an access descriptor are defined as follows:

Read Rights (bit 0). This bit indicates that reading the contents of the object referenced by this access descriptor is allowed.

Write Rights (bit 1). This bit indicates that writing the contents of the object referenced by this access descriptor is allowed.

Type Rights (bits 2-4). The interpretation of this 3-bit field is determined by the object type of the referenced object.

<u>Name of the bit</u>	<u>bit position in the AD</u>
Type Rights 1	bit 2
Type Rights 2	bit 3
Type Rights 3	bit 4

Local (bit 5). This bit indicates the object's lifetime. This bit is 0 for a global object and 1 for a local object.

Object Index (bits 6-31). This 26-bit field selects an object table entry in the object table.

Tag (Tag Bit). This bit must be 1 for a valid access descriptor.

Mixed Words

A mixed word is the union of a data word and an access descriptor, i.e. a mixed word can be viewed as either a data word or an access descriptor depending on the context.

The values of a mixed word are divided into the following classes:

(1) Valid Access Descriptor. A valid access descriptor has the tag bit set to 1. This can be dereferenced (used to reference the content of the object) if the object for the corresponding index is defined. The CONSTRAINT.INVALID_AD fault is raised when the tag bit is 0 when an AD is expected.

(2) Data. A data word has the tag bit set to zero. When a data value is generated, the tag bit is always set to zero.

When a data value is expected, the tag bit is ignored and interpreted as zero. No fault is raised when the tag bit is 1 when a data value is expected.

Virtual Addresses

FIGURE 3 is a diagram of a virtual address format. The fields of a virtual address are defined as follows:

Object Offset (bytes 0-3). This 32-bit field contains an Ordinal offset into the object referenced by the access descriptor in the virtual address.

Access Descriptor (bytes 4-7). This AD specifies the object referenced by this virtual address. The AD also specifies the permissible operations using this virtual address.

Object Table Entries

An object table (42) can contain the following types of object table entries. All object table entries are 16 bytes in size. Specific object table entries are identified by the entry type field (bits 96-98) of each object table entry as follows:

- 000 -- Invalid Object Table Entry
- 001 -- Embedded Descriptor
- 010 -- Invalid Simple Object Descriptor
- 011 -- Simple Object Descriptor
- 100 -- Invalid Paged Object Descriptor
- 101 -- Paged Object Descriptor
- 110 -- Invalid Bipaged Object Descriptor
- 111 -- Bipaged Object Descriptor

The last six object table entries are collectively called storage descriptors. The format of a storage descriptor is

shown in FIGURE 4. Valid storage descriptors contain physical addresses. Invalid storage descriptors, where the base address field may not be valid, are used to indicate that the selected object cannot be accessed.

The fields of a invalid simple object descriptor, a invalid paged object descriptor, or a invalid bipaged object descriptor are defined as follows:

TDO AD (bits 32-63). This field (51) contains the type definition object AD associated with this object descriptor.

Reserved (bits 68-69). This field must be zero.

Base Address (bits 70-95). This 26-bit field (52) contains the physical base address (in units of 64 bytes) of the block, page table or page table directory. This provides a 2^{32} byte physical address space. This field is uninterpreted in an invalid storage descriptor.

Entry Type (bits 96-98). This 3-bit field indicates the type of object table entries and the definition of the rest of the descriptor.

Access Status (bits 99-103). This 5-bit field is described in the next section. This field is only defined in a simple object descriptor. This field is preserved for other entry types.

Object Length (bits 114-119). This field contains the (encoded) length in units of 64 bytes of the storage block referenced by the base address field.

The word "encoded" is used with a length field to indicate that the field contains the actual length minus one instead of the actual length.

In a simple object, this field contains the (encoded) length in units of 64 bytes defined by this descriptor.

In a paged object descriptor, this field contains the (encoded) length in units of 64K bytes, or the number of page table entries in the page table (minus one)/16 defined by this descriptor.

In a bipaged object descriptor, this field contains the (encoded) length in units of 64M bytes, or the number of page table entries in the page table directory (minus one)/16 defined by this descriptor.

Write Back (bit 122). This bit enables write-back caching for the object if the segment or page is cacheable.

Object Type (bits 124-127). This 4-bit field encodes the object type of the object.

Access Status

The access status field contains information for the management of blocks and pages. It is found in simple object descriptors and valid page table entries (PTEs). This field does not appear in an invalid object descriptor, a paged/bipaged object descriptor, nor a page table directory entry (PTDE).

The 5 bits of the access status field shown in FIGURE 4 are defined as follows:

Accessed (bit 99 in OTE, Bit 3 in PTE). This bit indicates the object or page defined with this descriptor has been referenced (read or write). This bit is ensured to be 1 before the associated storage is referenced. This bit is

sticky (i.e. it is set implicitly by the system but never cleared).

Altered (bit 100 in OTE, Bit 4 in PTE). This bit indicates the object or page defined with this descriptor has been overwritten. This bit is ensured to be 1 before the associated storage is overwritten. This bit is sticky (i.e. it is set implicitly by the system but never cleared).

Mixed (bit 101 in OTE, bit 5 in PTE or PTDE). This bit indicates that an AD has been written in the object or page defined by this descriptor. This bit is ensured to be 1 before the associated storage is overwritten with a non-zero tag bit. This bit is sticky (i.e. it is set implicitly by the system but never cleared).

Cacheable (bit 102 in OTE, Bit 6 in PTE). This bit indicates the object or page defined with this descriptor can be cached. The encodings of the cacheable bit are as follows:

- 0 - Do Not Cache
- 1 - Can Be Cached.

Local (bit 103 in OTE, bit 7 in PTE). This bit indicates the lifetime of the object or page defined by this descriptor. This is 0 for a global object or page and 1 for a local object or page.

It may be necessary to set the accessed, altered, or mixed bits of the OTE/PTE in memory to ones on a memory access. The operation is performed using an atomic operation. The cacheability and local bits are never modified. It is permissible to alter these bits even though the OTE/PTE returned by the atomic read is invalid as long as the OTE/PTE is valid on a previous read and a TLB invalidation operation has not been performed on the specific OTE/PTE.

When an OTE/PTE is invalidated by software, only the valid bit should be modified first because the accessed, altered and mixed bits may be modified at the same time.

When the accessed, altered and mixed bits are updated in memory using RMW operations, the valid bit of the original OTE/PTE must not be altered.

Embedded Descriptors

Embedded Object Descriptor

An embedded descriptor holds special system-defined data structures. The only such data structure defined by the system is a semaphore.

The fields of an embedded descriptor are defined as follows:

Storage Area (bits 0-95). This 12-byte area is used to represent the data structure.

Entry Type (bits 96-98). This field is 2#001# for an embedded descriptor.

Use Default TDO (bit 99). If set, this descriptor has an associated default TDO. If clear, the TDO AD is assumed to be in bits 32-63 of the descriptor. (software would normally set this bit in an embedded descriptor.)

Type (bits 124-127). The only value defined by the system is 0100, which indicates that the first three words hold a semaphore.

Invalid Object Table Entry

The fields of an invalid object table entry are defined as follows:

TDO AD (bits 32-63). This field has the same interpretation as in a storage descriptor, but only if flag use-default-TDO is clear.

Entry Type (bits 96-98). This field is 2#000# for an invalid object table entry.

Use Default TDO (bit 99). If set, this descriptor has an associated default TDO. If clear, the TDO AD is assumed to be in bits 32-63 of the descriptor.

Object Type (bits 124-127). This field has the same interpretation as in a storage descriptor.

Page Table Entries

A page table or page table directory contains an array of 4-byte page table directory entries of similar format. Page table entries in a page table directory specify page tables while page table entries in a page table specify pages.

The fields of a valid page table entry (PTE) or page table directory entry (PTDE) are defined as follows:

Valid (bit 0). This bit is 1 to indicate a valid page table entry or page table directory entry.

Page Rights (bits 1-2). This 2-bit field encodes the permissible operations (read or write) in different execution modes on the contents of this page (in a page table entry) or for the pages by this page table (in a page table direc-

tory entry). Since a page may be controlled by more than one set of page rights, the effective rights is minimum of all page rights.

Access Status (bits 3-7). This 5-bit field is similar to that in a storage descriptor. This field is defined for a page table entry and is preserved for a page table directory entry.

Base Address (bits 12-31). This 20-bit field contains the physical base address (in units pages of 4096 bytes) of the page.

The field of an invalid page table directory entry is defined as follows:

Valid (bit 0). This bit is 0 to indicates an invalid page table directory entry.

Virtual Address Translation

For a memory request, an instruction specifies an Access Descriptor, an Object Offset, whether it is a Read or a Write access, and the Length of the Request.

The following describes the address translation of a virtual address to a physical address for accessing a paged object:

- (1) Compute the last byte of the memory request by adding the request length to the object offset.
- (2) Perform the request as two separate memory requests.
- (3) Determine the rep rights needed by the request type.

(4) Raise a PROTECTION fault if the rights needed are not presented in the read and write rights of the AD.

(5) Read the object table entry selected by the object index of the access descriptor. The entry type and valid bit are found in the encoded fields of the object table entry. Raise a DESCRIPTOR fault if the entry type is not appropriate. Raise a VIRTUAL MEMORY fault if the object table entry is not a valid storage descriptor.

(6) For this description it is assumed that the object table entry is for a paged object. Raise a PROTECTION fault if the directory index and page index found in the offset portion of the address of the memory request is greater than the object length in the paged object descriptor found in the page table.

(7) Scale the page index (bits 12-21 of the offset in the address) by 4 and add it to the base address in the object descriptor (in the object table) to form the physical address of the selected data page table entry (in the page table).

(8) Read the page table entry and raise a VIRTUAL_MEMORY fault if the page table entry is marked as invalid.

(9) Raise a RIGHTS PROTECTION fault if the access descriptor rights are greater than the page rights.

(10) Combine the page offset (bits 0-11 of the address offset field) with the base address in the page table entry in the page table to form the physical address of the paged object referenced by the memory request.

Object Referenc Testing

COMPARE_MIXED and CHECK_TAG

The COMPARE_MIXED instruction compares both ADs or data for equality. If both operands are ADs, the instruction tests whether they reference the same object. If both operands are data, the instruction test whether the data value are equal. The CHECK_TAG instruction checks for the tag bit.

In an untagged system in supervisor mode, the tag bit is assumed to be set in COMPARE_MIXED.

In the Ada expression "access_type_variable = null" where AD is used to represent access variable, the CHECK_TAG should be used instead of the COMPARE_MIXED with zero because any non-zero data values still cannot be used to reference an object.

Access Descriptor Creation

The protection system requires that the tag bit not be set to one arbitrarily, that lifetime checking not be bypassed, and that the mutator operation not be ignored. However, in certain controlled situations, system software needs to create an access descriptor (AD) for an existing object. This is accomplished by the CREATE_ACCESS_DESCRIPTOR instruction.

The CREATE_ACCESS_DESCRIPTOR instruction converts a data word to an AD.

Object Addressing Instructions

LOAD_PHYSICAL_ADDRESS

The LOAD_PHYSICAL_ADDRESS instruction returns the physical address of the operand.

Type Management and Access Control

Typed Objects

The present invention provides two typing mechanisms: system-type and extended-type. They are not mutually exclusive. System types can be viewed as predefined properties of the object, while the extended type provides a unique identifier of the type and allows for type-specific software defined functions. The system-type mechanism is to provide system type specific instructions, which require operands of specific system types. The extended-type mechanism supports user defined types. This allows software defined type specific operations and extended-typed specific operations to verify an object's type before carrying out their prescribed functions. This facility supports the type-manager style of programming.

The type definition object (74) normally contains a reference (75) to a type manager domain object (71) that provides some type-specific procedures.

Object Type Field

The system type of an object is specified by the object type field in its object descriptor (see application SN 07/455,625 (D-1256)). One of the system types is generic (i.e. a simple flat address space) which contains no architecture defined fields.

The encodings for the Object Type field are as follows:

<u>Encoding</u> (Binary)	<u>Object Type</u>
0000	Generic
0001	Type Definition Object
0010	Process Object
0011	Domain Object
0100	reserved
0101	Port Object
0110-0111	reserved
1000-1111	available for system software

The architecture defines one type value for an embedded descriptor: 0100, which denotes a semaphore.

Type Definition Object (TDO)

Refer to FIGURE 1. The extended type of an object (38) is specified by the type definition object's access descriptor (51) associated with the object descriptor in the object table (42), either explicitly or via default. If a default type definition object (TDO) is to be used, the object index for the default TDO is 16 plus the type encoding of the object or embedded descriptor. For instance, for a semaphore, the default TDO is entry 20 in the object table.

A type definition object contains a reference address (75) to the type manager domain (71). The type manager domain (71) contains other information used to manage the instances of a particular object type. These type manager domains reference software defined, type specific operations, that are programmed as procedures (77) or, in some cases, as independent processes.

The object index of the TDO AD (51) can also be used as an unique identifier for the extended type. In this manner, up to 2^{26} (64 mega-) extended types can be defined.

For system objects, while the object type is used to indicate the set of type specific instructions, the type definition object can be used to provide software (i.e. procedural) extensions to the architecture-defined, type specific operations.

Different type definition objects may be associated with different instances of the same system object type. This permits software extensions to the architecture defined instructions to be defined on a per instance basis.

LOAD_TYPE_DEFINITION

This instruction copies the type definition object AD associated with the object referenced by the source AD into the destination. If the object specified has a default TDO, an AD (with no type nor rep rights and specifying global lifetime) to the default TDO is returned.

Rights

The architecture uses various rights bits to restrict the way in which an object or an AD may be manipulated. There are two sets of rights bits -- one set is found in an AD and the other set is located in the various page table entries in the access (i.e. address translation) path.

Different ADs can have different access rights to the same object. Rights on an access path are shared among all users of the access path. An AD can be easily duplicated with the same or less rights than the source AD. Since an access path cannot be similarly duplicated, the architecture makes any

change to rights in the access path affect all users of the access path.

An AD contains the following rights bits:

- Type Rights
- Read Rights
- Write Rights

A page table entry in an access path contains the following rights bits:

- Page Rights

Type Rights

The type rights of an AD define the permissible type-specific operations for the object referenced by the AD. The interpretation of the type rights bits of an AD depends on the type of object referenced. The type rights bits for a generic object are uninterpreted. The interpretation of the type rights bits for each system object are predefined. They are described in the individual system object descriptions throughout this specification. The uninterpreted type rights bits are preserved for software defined type rights and interpreted by individual software level type managers.

Read and Write Rights in Access Descriptors

The read rights control reading the content of the referenced object, while the write rights control writing. The actual permissible operations are also optionally determined by the page rights in page tables. This allows the support of inaccessible objects (i.e. no access for encapsulated objects), read-only objects, write-only objects, and read-writeable objects.

Execute rights are not defined; only read rights are required to execute an instruction in the execution environment.

Page Rights

The page rights in a valid page table entry define the read/write rights of the page in a bipaged object or a paged object. Page rights are used to permit software defined areas or regions of the execution environment to have different access protection.

Page rights are interpreted differently depending on the execution mode of the current process and the PTE break mode of the processor. Page rights are defined as follows for the supervisor and user execution modes:

PTE Break	User Mode	Supervisor Mode
1	00 - read-write	read-write (break on write)
0	00 - no access	read-only
X	01 - no access	read-write
X	10 - read-only	read-write
X	11 - read-write	read-write

Page rights, instead of AD read/write rights, are used to protect the instruction areas of a linear address space from accidental modification.

When the processor is in physical-addressing mode (address translation is disabled), rights checking is disabled.

As described above, the access rights for an item are defined by rights fields in each level of the access path. Each access path contains the following:

Read and write rights in an AD.

Page rights in the page table directory entry for a bipaged object.

Page rights in the page table entry for a paged object or a bipaged object.

The effective rights for an object address are the minimum of the rights in the access path.

INSPECT_ACCESS

The INSPECT_ACCESS instruction returns the effective page rights of the access path specified by the source address.

Type Definition Object

A type definition object (74) is used to control access rights amplification. A type definition object has a predefined system type.

The type rights in an AD (51) for a type definition object are defined as follows:

Type Right 1 Uninterpreted.

Type Right 2 Amplify Rights: If the bit is 1, the TDO may be used in the AMPLIFY_RIGHTS instruction.

Type Right 3 Create Rights: If the bit is 1, the TDO may be used in the CREATE_ACCESS_DESCRIPTOR

instruction.

The fields of the type definition object (74) shown in FIGURE 1 are defined as follows:

Super TDO (bit 0). This bit is interpreted during rights amplification as follows:

0 - This TDO can be used to amplify ADs for objects whose type matches that specified by this TDO.

1 - This TDO can be used to amplify any AD.

Extended (bit 1). This bit is 1 if the TDO is used to manage objects having this object as their TDO ADs. This bit is 0 if the TDO is used to manage objects having the same object type as specified in the TDO.

This bit also allows a single TDO, to manage object of the specified object type independent of their TDO ADs.

Object Type (bits 28 - 31). This 4-bit field is defined to have the same format as the corresponding field in an object descriptor. In the AMPLIFY_RIGHTS instruction, this field is compared against the object type of the referenced object if the extended type and the super TDO bits are zero.

Rights Manipulation

Refer to FIGURE 8. Two instructions are defined to increase or decrease the access rights of an AD. A rights mask (170) is specified in the source_2 operand of the instruction during rights amplification and restriction. A rights mask has the same format as an access descriptor, except that the lo-

cal bit, the object index field and the tag bit are not used.

If the `AMPLIFY_RIGHTS` instruction (165) is decoded, access rights are amplified by logically-ORing, in the amplify rights logic (168), the access rights in the source AD (162) with those in the rights mask (170). The logic flow within the amplify rights logic (168) of FIGURE 8 is shown in FIGURE 9.

If the `RESTRICT_RIGHTS` instruction (167) is decoded, , access rights are reduced by logically-ANDing, in the restrict rights logic (172), the access rights in the AD with the complement of that in the rights mask (170). The logic flow within the restrict rights logic (172) of FIGURE 8 is shown in FIGURE 10.

The `AMPLIFY_RIGHTS` instruction requires a type definition object AD (51) with amplify rights. The `AMPLIFY_RIGHTS` instruction amplifies the rights of the source AD (162) as specified by the rights mask (170) and stores the amplified AD in the destination (176).

The `RESTRICT_RIGHTS` instruction removes the rights of the source AD (174) as specified by the rights mask (170) and stores the restricted AD in the destination (176).

The accessibility of an addressing environment can be restricted to those objects with the minimum required access rights needed for the execution of the program. Instructions are provided to remove (or restrict) access rights that are not needed. Access restriction is typically performed before an AD is passed as a parameter to a procedure or returned as a result.

Certain (system or xtended) typed objects are manipulated exclusively by their corresponding type managers. An AD to a typed object, outside the domain of its type manager, usually has limited access rights (e.g. no access or read-only rights) to prevent access to or modification of the object without the knowledge of the type manager. The type manager restricts (i.e removes) the access rights in the ADs, for objects of the type(s) under its control, before passing them outside its domain. When such an AD is returned to the type manager, the access rights are amplified (i.e. increased) while inside the domain to allow modification, of the referenced object, by the type manager.

AMPLIFY_RIGHTS instruction

Refer to FIGURE 9 which is a diagram of data flow within the amplify rights logic of FIGURE 8 for an understanding of the following description of the AMPLIFY_RIGHTS instruction.

Operation: Amplifies, under control of a type definition object, the specified rights bits in the specified source AD and stores the amplified Ad in the destination.

Opcodes:

amplify

653 REG	source_1(val:m)	source_2(val:0)	src/dst	(vds:m)
	source_1		type definition object	
	source_2		rights_mask	
	src/dest		source/destination	

Action:

Check the tag bit of the TDO AD (202). Raise a CONSTRAINT_INVALID AD if the tag bit of the source AD is zero (204).

Check the typ rights field to see if amplify rights are specified (206). If not, assert a PROTECTION TYPE RIGHTS fault (208).

Check the super TDO bit (210). If the super TDO bit in the TDO is one, go to logic block (218). If the super TDO bit in the TDO is zero and the extended bit is one (212), raise a TYPE.MISMATCH FAULT (216).

Compare the object type of the source object with the object type controlled by the TDO (214). If different, and the super TDO bit in the TDO is zero (21) and the extended bit is zero (212), then raise a TYPE.MISMATCH FAULT (216).

If the object type of the source object with the object type controlled by the TDO (214) are the same, then Logically OR the read, write and type bits of the source AD and the corresponding bits in the rights mask to form the type rights of the destination AD (218).

RESTRICT_RIGHTS instruction

Refer to FIGURE 10 which is a diagram of data flow within the restrict rights logic of FIGURE 8 for an understanding of the following description of the RESTRICT_RIGHTS instruction.

Operation: Restricts the specified rights bits in the specified source AD and stores the Restricted AD in the destination.

Opcodes:

restrict

652 REG source_2(val:0) src/dst (vds:m)

source_2 rights_mask

src/dest source/destination

Action:

Check the type rights field to see if restrict rights are specified (252). If not, assert a PROTECTION TYPE RIGHTS fault (254).

If the tag bit of the source is zero (256) skip the following rights restriction actions and store the data segment in the destination (260).

If the tag bit of the source is one (256) remove the read, write and type bits of the source AD whose corresponding rights bits in the rights mask are one (260). Store the restricted AD in the destination (260).

Cacheable

The cacheable bit is defined to prevent "caching" of an object or a page. This bit is one if caching is allowed, zero otherwise. Caching allows delaying a memory access (e.g. write back), and grouping contiguous memory accesses into a single memory transaction (e.g. cache fill). If caching is disabled, memory accesses are passed directly to the physical memory without any modification.

In a system with memory mapped device registers, the cacheable bit for the page or object should be zero. It may also be used to prevent external caching of instruction region to avoid double caching, or to prevent caching of shared objects to avoid unnecessary caching traffic.

The cacheable bit in a simple object descriptor or a page table entry indicates whether the contents of the referenced object or page can be cached.

While the invention has been particularly shown and described with reference to preferred embodiments thereof, it will be understood by those skilled in the art that the foregoing and other changes in form and detail may be made therein without departing from the scope of the invention.

CLAIMS

1. A data processing system including a memory space (10) and a processor register space (11) in a processor of said data processing system for storing object descriptors having base and length information stored therein for use in locating objects in said memory space,

said processor being capable of communicating with said memory space (10) and capable of executing an operation by means of an operator specified in an instruction (14);

a plurality of addressable objects, including a bipaged object (38), said objects being stored in an address space of said memory space,

said plurality of addressable objects including instruction objects comprised of instructions including said instruction (14) defining an operation;

an address (20) referenced in said instruction, said address being specified as an offset (22) and an access descriptor (24);

said plurality of objects including an object table (42), a page table directory (60), and a page table (44);

each of said objects being comprised of words, said words being either access descriptors (30) or data words (32);

each of said access descriptors including

1. an object index for selecting an object in said address space,

2. a rights field, said rights field specifying the permissible operations on an object selected by said access descriptor;

said data processing system characterized by:

each of said access descriptors including an AD lifetime bit for indicating the lifetime of an object specified by said access descriptor;

said object table object (42) being stored at an address accessible by said processor upon processor initialization,

said object table having stored therein object descriptors for use by said processor in forming physical addresses to said page table directory object (60),

each of said object descriptors including a first encoded field (50) specifying the base address of the object referenced by said object descriptor and a second encoded field (52) which includes page rights;

said object storage descriptor includes a lifetime bit for indicating the lifetime of said object selected by said access descriptor;

said page table (44) having stored therein page table entries for use by said processor in forming physical addresses to said paged object (38),

each page table entry having a page rights field (81);

said offset (22) including a page index (PI) used to index into a selected page table to locate said page table entry, and a page offset (PO) used in conjunction with a base address (79) stored in said page table entry to locate a datum entry in said paged object (38); and,

means (150) for comparing said AD lifetime bit of said access descriptor with the said OTE lifetime bit of said object table entry and for asserting a fault upon the condition that the access permitted by said AD lifetime bit of said access descriptor is inconsistent with said OTE lifetime bit of said object table entry.

2. The combination in accordance with claim 1 wherein said AD lifetime bit for indicating the lifetime of an object specified by said access descriptor comprises a

Local bit (bit 5), said local bit having a value of 0 indicates a global object with unbound object lifetime and said local bit having a value of 1 indicates a local object with bound object lifetime.

3. The combination in accordance with claim 2 wherein wherein said means (150) for comparing said AD lifetime bit of said access descriptor with the said OTE lifetime bit of said object table entry includes means for asserting a PROTECTION.LIFETIME fault upon the condition that said AD has its associated local bit set to 1 and the lifetime bit of said object table entry (OTE) has its associated local bit set to 0 (is of global lifetime).

4. A data processing system substantially as hereinbefore described with reference to the accompanying drawings.